

FACTORING OF N=pq USING k

In several earlier notes we have found that a semi-prime $N=pq$, with primes $p<q$, can be factored into the forms $p=6n\pm 1$ and $q=6m\pm 1$, where n and m are integers whose values depend on finding integer solutions of –

$$R[k] = \sqrt{(H + 6k)^2 - 4(B - k)} \quad \text{when } N \bmod(6) = 1$$

or

$$S[k] = \sqrt{(H + 6k)^2 + 4(B - k)} \quad \text{when } N \bmod(6) = 5$$

Here $A=(N-1)/6$ for the first case and $A=(N+1)/6$ for the second case. Also $H=A \bmod(6)$ and $B=(A-H)/6$. For smaller and intermediate sized N s the above radicals are easy to solve to produce integer values. However, when N is large it becomes difficult to find the right value of variable k which allows this. The values of n and m are given as-

$$[n, m] = \frac{1}{2} [(H + 6k) \pm R] \quad \text{or} \quad [n, -m] = \frac{1}{2} [(H + 6k) \pm S]$$

For a typical semi-prime the quantities B and H will be known, so one needs to only find the value of k which makes the radical an integer. Although B is typically much smaller than k , k can nevertheless become large so that a brute force evaluation of one or the other of the radicals can become extremely time consuming.

We show here how to get around this difficulty by estimating a value for k designated by k_1 . The procedure works as follows. It is known that –

$$p = 6n \pm 1 = \alpha \sqrt{N} \quad \text{and} \quad q = 6m \pm 1 = (1/\alpha) \sqrt{N} \quad \text{with} \quad 0 < \alpha < 1$$

Thus for large N , we can say that-

$$n \approx \alpha \sqrt{N}/6, \quad m \approx (1/\alpha) \sqrt{N}/6, \quad \text{and} \quad p/q \approx \alpha^2$$

The range for α is $0 < \alpha < 1$ with $\alpha=1$ meaning that $n=m$ and $p=q$.

Now we can get an estimate for the desired value of k by eliminating n from its two definitions. For the case of $N \bmod(6)=1$, we get-

$$\frac{\alpha \sqrt{N}}{6} = \left(\frac{1}{2} \right) \{ (H + 6k) - R \}$$

and for $N \bmod(6)=5$, we have-

$$\frac{\alpha\sqrt{N}}{6} = \left(\frac{1}{2}\right)\{(H + 6k) - S\}$$

Solving for k in these last two expressions, we have, after noting $H \ll 6k$ and $1 \ll \alpha\sqrt{N}$, that -

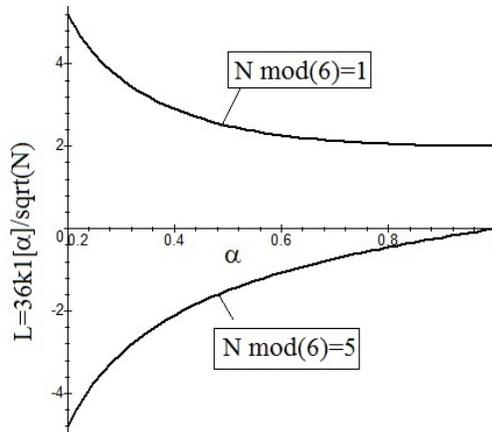
$$k1[\alpha] \approx \frac{\alpha\sqrt{N}}{36} \left\{1 + \frac{1}{\alpha^2}\right\} \quad \text{for } N \bmod(6) = 1$$

and

$$k1[\alpha] \approx \frac{\alpha\sqrt{N}}{36} \left\{1 - \frac{1}{\alpha^2}\right\} \quad \text{for } N \bmod(6) = 5$$

These values can now be used to search for the k which should lie close to k1 and which makes the radical an integer. The values for k1 depend not only on the root of N but also on α . Typically large semi-primes will require large k1s and the size of |k1| will increase with decreasing α . The following graph characterizes this behavior-

NON-DIMENSIONAL K1 VERSUS ALPHA



In the graph we have run the non-dimensional quantity $L=36k1[\alpha]/\sqrt{N}$ over the range $0.2 < \alpha < 1$ for both types of semi-primes. The increase in L with decreasing α is at first small but then increases rapidly for values of α in the given range. Typically we have $L \approx 2$ for $N \bmod(6)=1$ and $L \approx -1$ for $N \bmod(6)=5$ provided α lies between 0.5 and 1. We also find the unique value of $L=2$ occurring for those semi-primes N where $p=q$.

Let us next demonstrate the above points by working out a few explicit factorizations of larger semi-primes. Take first -

$N=155505643$ where $\sqrt{N}=12470.19$, $N \bmod(6)=1$, $A=(N-1)/6=25917607$, $H=A \bmod(6)=1$, and $B=(A-H)/6=4319601$.

We assume first that $\alpha=1$, so that we have the $k1[1]$ estimate -

$$k1[1] = \sqrt{N} / 18 = 692.79$$

If $\alpha < 1$, then $k1[\alpha]$ increases to a value of $\sqrt{N}/36 \{ \alpha + 1/\alpha \}$.

Now carrying out the following computer search-

for k from 693 to 723 do {k,sqrt((1+6*k)^2-4*(4319601-k))}od

yields $R=1017$ at $k=713$. So we have our solution -

$$[n, m] = (1/2)(1+6(713) \pm 1017) = [1631, 2648]$$

which means that-

$$155505643 = \{6(1631)+1\} \{6(2648)+1\} = 9787 \times 15889$$

As a side benefit, we now know the value of α . It equals $\alpha = \sqrt{p/q} = 0.7848$.

If we had used this value of α instead of $\alpha=1$ then $k1[0.7848]=713.22$ and so essentially matches the solution of $k=713$. The advantage of using $k1[1]$ in our calculations is that we know for the $N \bmod(6)=1$ case that it offers a lower bound on the actual $k1[\alpha]$ considerably larger than $k=0$.

Another interesting point following from the $N \bmod(6)=1$ case is that $k1[1]=\sqrt{N}/18$ for all positive integer semi-primes including the one hundred digit long N s used in public key cryptography. So for a semi-prime of 100 digit length the k to be used in an integer search for R must be some 48 digit long or longer $k1[1]$.

Consider next an N where $N \bmod(6)=5$. Such an example is-

$$N=3475379339=(6n-1)(6m+1) \text{ where } N \bmod(6)=5 \text{ and } \sqrt{N}=58952.348.$$

Here $A=(N+1)/6=579229890$, $H=0$ and $B=(A+h)/6=96538315$. Although the $k1[1]$ case is here equal to zero, one can use the neighboring value corresponding to $\alpha=0.8$. This yields the search starting point of $k1[0.8]=-736.90$. Carrying out a search with this last value leads to an integer $S=20314$ at $k=-858$. That is it took 122 operations to find an integer answer. the rest of the problem is now easy. We have-

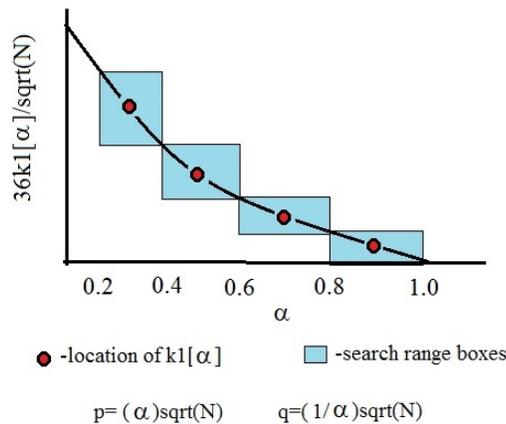
$$[n, -m] = (1/2) \{ (-6(858) \pm 20314) \} = [7583, 12731]$$

From this result follows the factorization-

$$3475379339=45497 \times 76387$$

In both of the above case the factoring of nine and ten digit long semi-primes was fairly easy to accomplish compared to other existing methods such as elliptic curve factorization and generalized grid methods. It is very likely that even larger semi-primes can be factored by the present approach. To prevent rapid factorization will require that N have values of α lying in a range $0 < \alpha < 0.1$. There it becomes difficult to find a k_1 close to k since α is not known before hand. The most consistent approach to factoring by the present method is to determine several different $k_1[\alpha]$ values for a given N and then search within a limited range for integer solutions of R or S for a given $k_1[\alpha]$. If no solution appears after a 10% search range proceed on to the next $k_1[\alpha]$. This stepping procedure will eventually lead to an integer solution for R or S and thus factorization. Here is a schematic of such a search method for the $N \pmod{6}=1$ case-

SEARCH RANGES FOR FINDING INTEGER R OR S



One can start the search for integer R or S within any of the boxes centered on a given $k_1[\alpha]$.

U.H.Kurzweg
 April 17, 2017