

A GRAPHICAL TECHNIQUE FOR FACTORING LARGE SEMI-PRIMES

One of the more important unsolved problems in number theory is how to quickly factor a large semi-prime $N=pq$ into its two component p and q . Our own efforts over the last decade concerning the factoring of such semi-primes has led to the closed form result-

$$[p,q]=a\pm\sqrt{a^2 - N}$$

, where $a=(p+q)/2=[\sigma(N)-N-1]/2=Nf(N)/2$. Here $\sigma(N)$ is the sigma function of number theory and $f(N)$ the number fraction discovered by us earlier. The definition of $f(N)$ is-

$$f(N)=\sigma(N)-N-1)/N$$

It is a slowly increasing function of N with $f(p)$ and $f(q)$ both equal to zero. Explicit values for p and q are thus obtainable if $\sigma(N)$ or $f(N)$ are known. One is fortunate in that most advanced math computer programs, such as Maple or Mathematica, give sigma for values of N as high as 40 digit length, meaning that prime components as high as twenty digits each can be found by the above $[p,q]$ formula. For still larger semi-primes, such as found in public key cryptography, some additional work on quickly finding sigmas for N s above forty digits length is needed.

It is the purpose of this note to introduce a new graphical approach for factoring large semi-primes. Hopefully this will offer some clues as to how to find larger $\sigma(N)$ for N s of greater than forty digit length.

Our starting point is the basic definition for any semi-prime-

$$N=pq$$

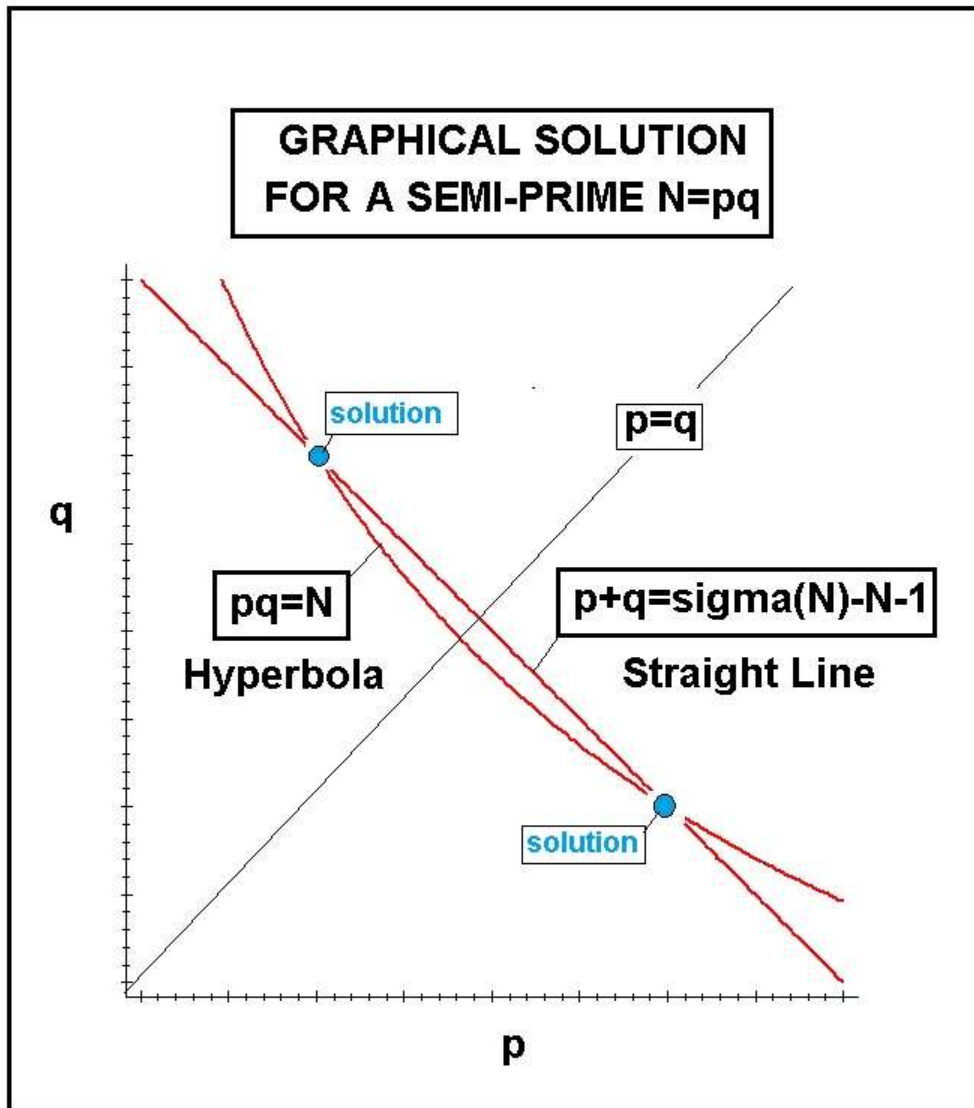
When looking at this function graphically we have a hyperbola symmetric about line $p=q$. Next we introduce the sigma function for N with the property that-

$$\sigma(N)=\sigma(p)\sigma(q)=(p+1)(q+1)=N+(p+q)+1$$

Expanding this definition we get-

$$(p+q)=\sigma(N)-N-1$$

Graphically this represents a straight line in the p - q plane which cuts the hyperbola at two points both representing the solution $[p,q]$. Here is the combined graph-



Note that the straight line is also symmetric about $p=q$.

A more convenient way to describe our solution is to eliminate q to get-

$$y=x^2-x(\sigma(N)-N-1)+N$$

in the x - y plane. Here $p=x$ and $y=0$ corresponds to our two solutions. Note that this $y=y(x)$ equation is a parabola which may be written as-

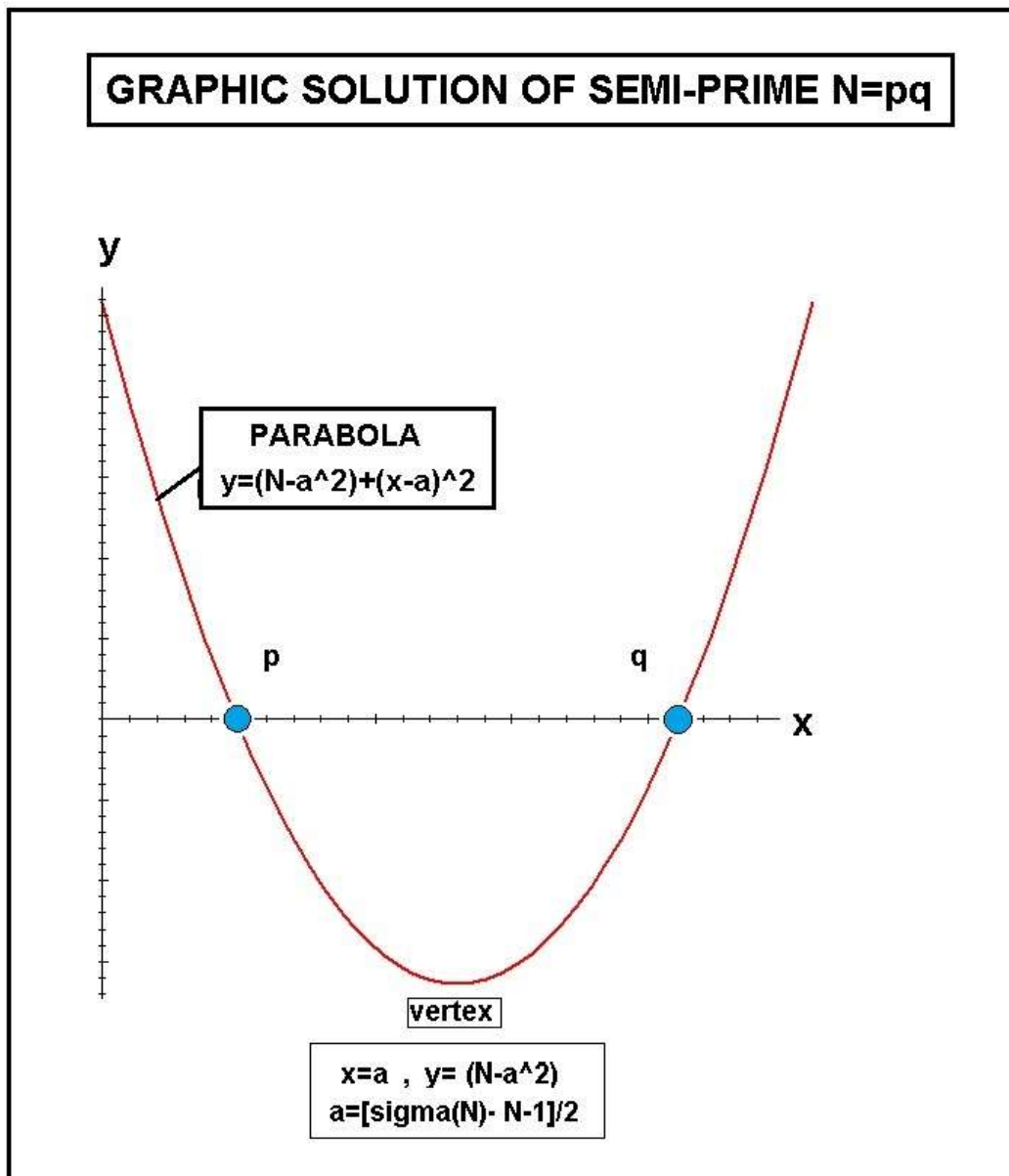
$$y=(N-a^2)+(x-a)^2$$

Its vertex lies at $x=a$ and $y=N-a^2$. Here-

$$a=[\sigma(N)-N-1]/2=Nf(N)/2$$

as already defined earlier.

This last expression is an important parameter which locates the parabola in the x-y plane. Here is a plot of the parabola for a typical semi-prime $N=pq$ -



The mean x value between the two solutions is $x=a$. Also note that 'a' is approximately equal to \sqrt{N} . Let us demonstrate this graphical approach for factoring a couple of specific large semi-primes.

The first example involves factoring the semi-prime-

$$N=455839 \text{ .}$$

With aid of our Maple program, we find $\sigma(N)=457200$ in a split second. This information then yields-

$$a=457200-455839-1=1360$$

This leads to the parabola-

$$y=(N-a^2)+(x-a)^2=-6561+(x-680)^2.$$

On setting y to zero we get the solution-

$$x=680 \pm 81$$

That is-

$$p=599 \text{ and } q=761$$

This particular N has been used in the literature to demonstrate the elliptic curve factorization method of Lenstra for semi-primes. It takes considerable more effort by that method to arrive at the same result for p and q.

As a second factorization consider the Fermat Number-

$$N=2^{32}+1=4294967297$$

which Fermat thought to be prime but proven later by Euler to be a composite.

Its sigma function value is given by our Maple program in a split second and reads-

$$\sigma(N) = 4301668356$$

This produces-

$$a=[\sigma(N)-N-1]/2=3350529$$

Next setting y to zero, we arrive at $x = a \pm \sqrt{a^2 - N}$. This in turn leads to the solution-

$$p = x = 641 \quad \text{and} \quad q = N/p = 6700417$$

What took Euler weeks to arrive at is here gotten in a split second.

Although we could factor numerous number of additional semi-primes for N s up to about 40 digit length, what is now needed is to find a way to speed up the search for $\sigma(N)$ s for $N > 10^{40}$. If this can be archived, then present day public key cryptography will become obsolete.

U.H.Kurzweg

July 22, 2024

Gainesville, Florida